



JNCF 2018, *Marseille*

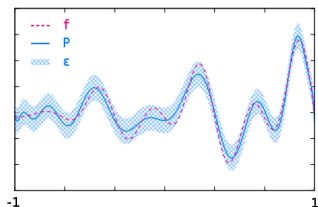
Multinorm Validation and Vector-Valued D-Finite Functions

F. Bréhard

Why Algorithmic and Rigorous Polynomial Approximations?



- ▶ Rigorous Polynomial Approximation =
Polynomial + error bound

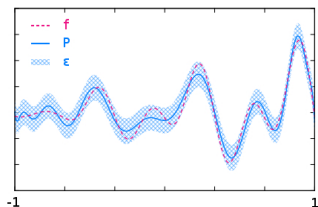


- Rigorous methods
- Algorithmic methods
- Efficient and accurate
- To be integrated in a large-scale library

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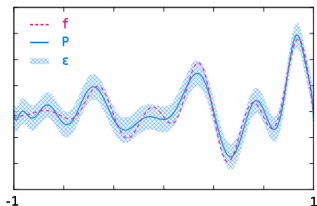
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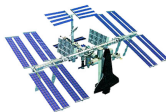
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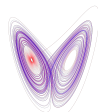
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- ▶ Solutions of **coupled** systems of linear ordinary differential equations.
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- ▶ Various fields of applications:



Safety-critical
engineering



Computer-aided
mathematics

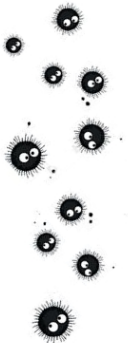
Outline

1 Introduction

2 Multinorm Validation: a New Framework

3 A Posteriori Validation of Vector-Valued D-Finite Functions

4 Conclusion and Future Work



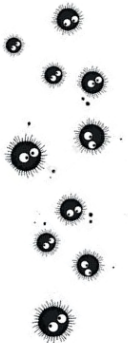
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General scheme

- ▶ Fixed-point equation $\mathbf{T} \cdot x = x$ with \mathbf{T} contracting,
- ▶ Approximation x to exact solution x^* ,
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Banach Fixed-Point Theorem

If (X, d) is complete and \mathbf{T} **contracting** of ratio $\mu < 1$,

- ▶ \mathbf{T} admits a unique fixed-point x^* , and
- ▶ For all $x \in X$,

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Sketch of the proof:

$$\blacksquare \quad d(x, x^*) \leq d(x, \mathbf{T} \cdot x) + d(\mathbf{T} \cdot x, x^*)$$



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Quasi-Newton Method for $\mathbf{F} \cdot x = 0$

Compute $\mathbf{A} \approx (\mathbf{DF})_x^{-1}$ in order to define:

$$\mathbf{T} \cdot x = x - \mathbf{A} \cdot \mathbf{F} \cdot x.$$

Banach fixed-point theorem applies if for some $r > 0$:

- $\mu = \sup_{\tilde{x} \in B(x, r)} \|\mathbf{1} - \mathbf{A} \cdot \mathbf{DF}_{\tilde{x}}\| < 1,$
- $\|x - \mathbf{T} \cdot x\| + \mu r < r.$



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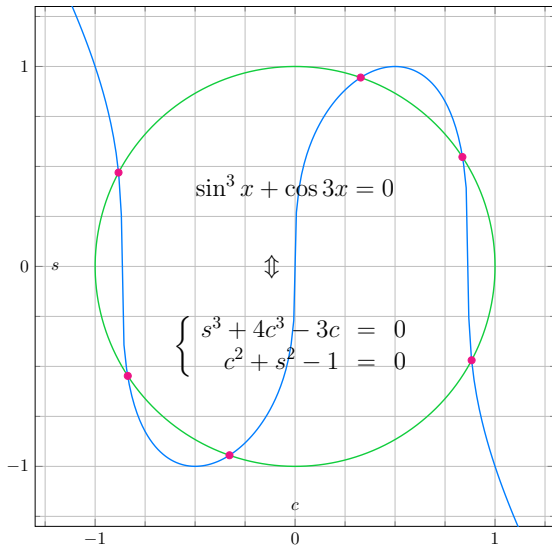
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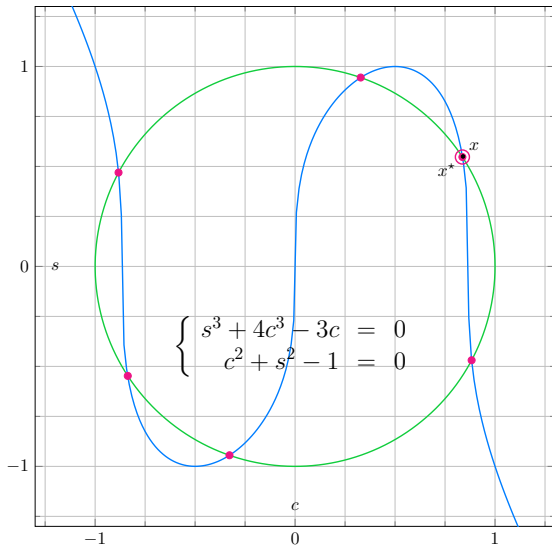
Applications to function space problems:

- Early works by Kaucher, Miranker, Yamamoto *et al* (~80's, ~90's).
- Lessard *et al* (2007 - today).
- Benoit, Joldes, Mezzarobba (2011)
 Bréhard, Brisebarre, Joldes (2017).

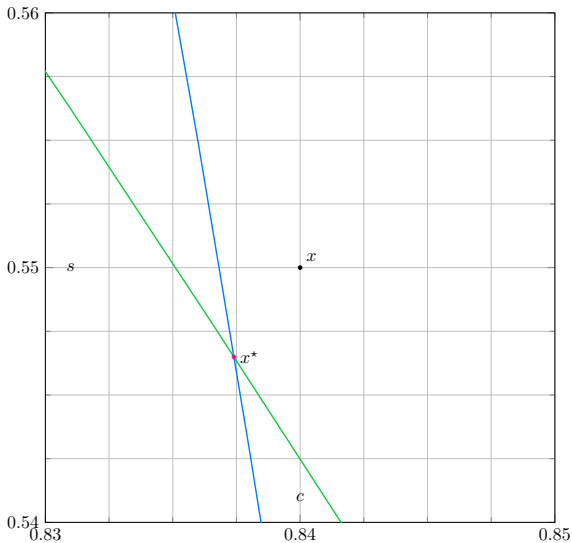
Example: Polynomial Equation in the Plane



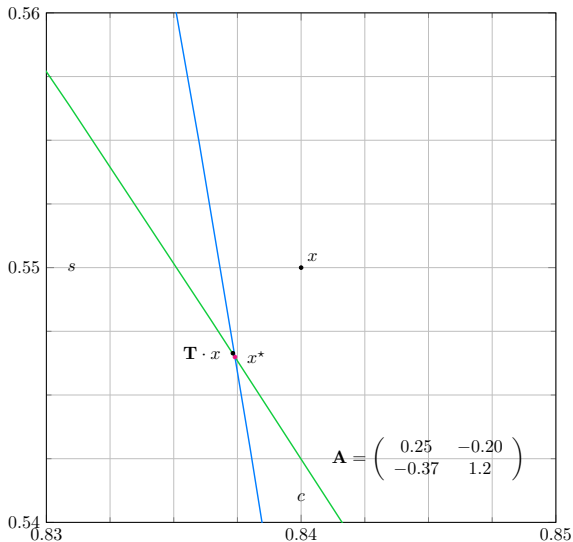
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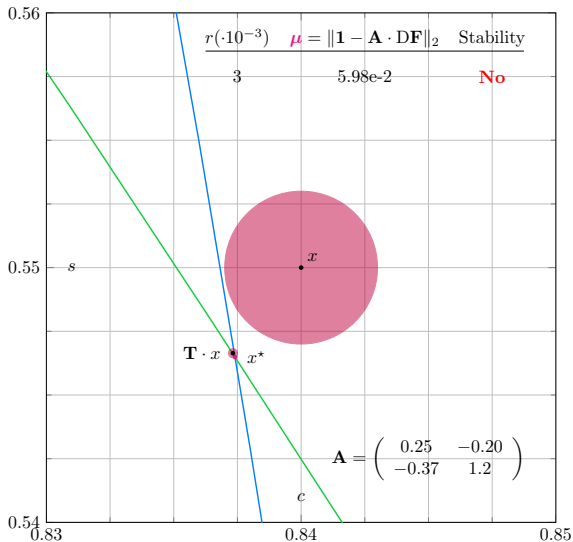
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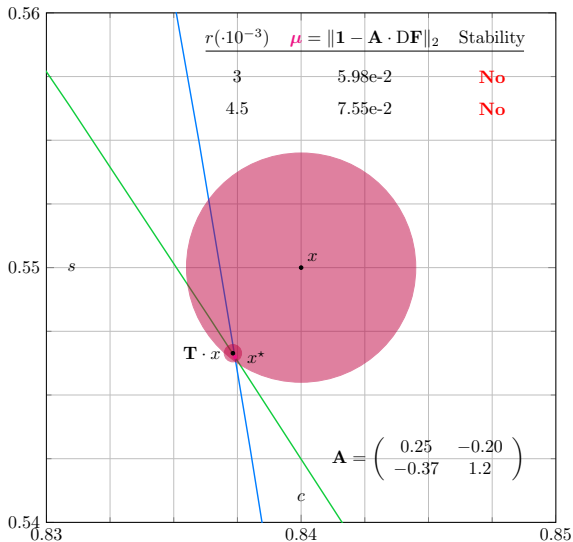
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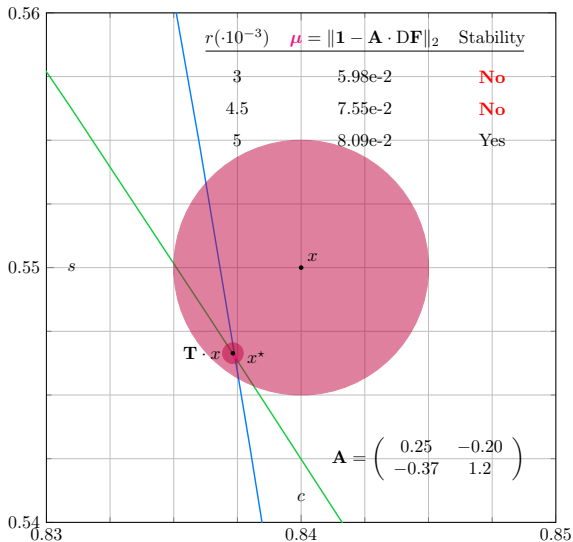
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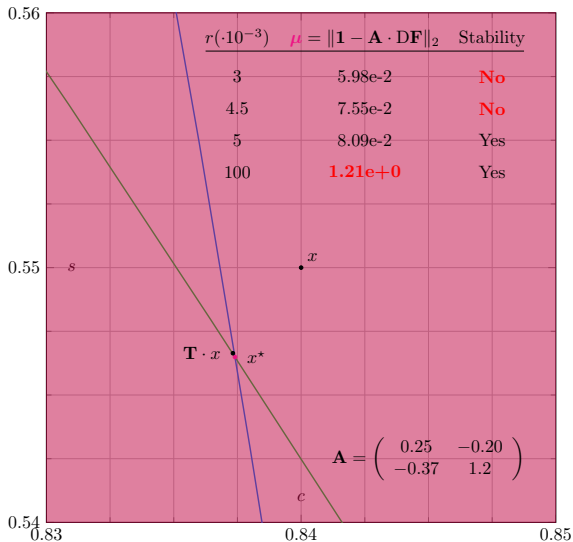
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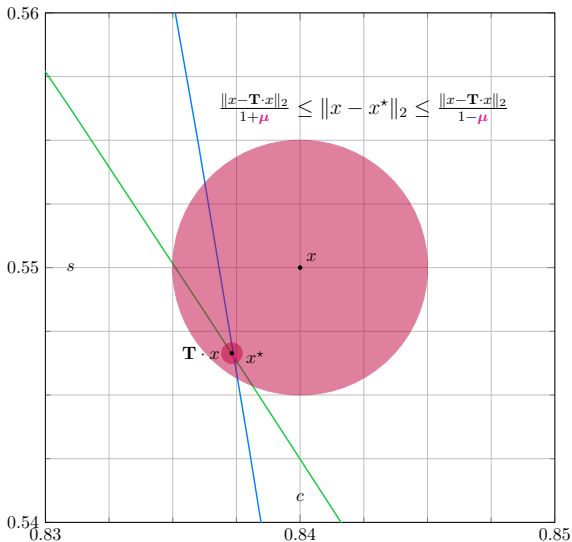
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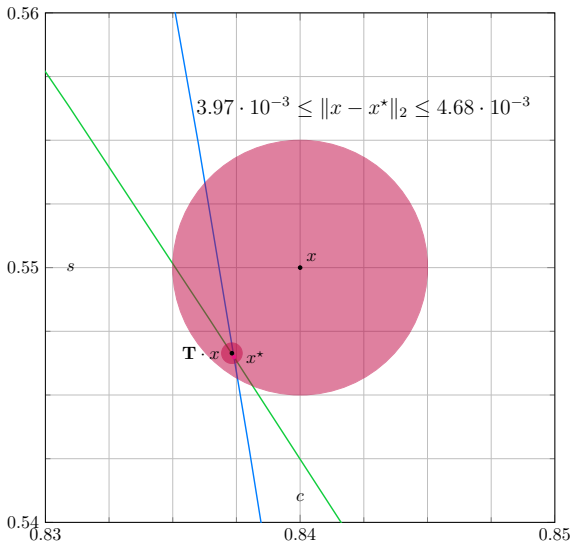
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Vector-Valued Metric

$(X_i, d_i)_{1 \leq i \leq p}$ complete metric spaces.

- $d(x, y) = (d_1(x_1, y_1), \dots, d_p(x_p, y_p))$
 $\in \mathbb{R}_+^p$ *vector-valued* metric.
- $\mathbf{F} : X \rightarrow X$ is $\mathbf{\Lambda}$ -Lipschitz for $\mathbf{\Lambda} \in \mathbb{R}_+^{p \times p}$:

$$d(\mathbf{F} \cdot x, \mathbf{F} \cdot y) \leq \mathbf{\Lambda} \cdot d(x, y) \quad \forall x, y \in X$$



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Convergent to Zero Matrices

$\Lambda \in \mathbb{R}^{p \times p}$ is convergent to zero if:

- $\Lambda^k \rightarrow 0$ as $k \rightarrow \infty$,
- $\Leftrightarrow \rho(\Lambda) < 1$.

Generalized Contractions

$f : X \rightarrow X$ is a *generalized contraction* if it is Λ -Lipschitz for Λ convergent to zero.



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$$\blacksquare (1 - \Lambda)^{-1} = 1 + \Lambda + \Lambda^2 + \dots + \Lambda^k + \dots \geq 0.$$

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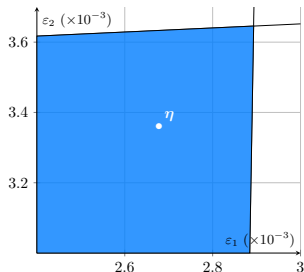
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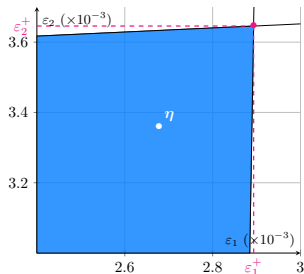
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$$(1 + \textcolor{red}{\Lambda}) \cdot d(x, x^*) \geq d(x, \mathbf{T} \cdot x)$$



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$$\blacksquare (1 + \Lambda)^{-1} = 1 - \Lambda + \Lambda^2 - \dots + (-1)^k \Lambda^k + \dots \not\geq 0.$$

⇒ Cannot deduce lower bounds!



Error Polytope

Let $\epsilon = d(x, x^*)$ and $\eta = d(x, \mathbf{T} \cdot x)$:

$$(1 - \Lambda) \cdot \epsilon \leq \eta \quad (\text{P})$$

$$(1 + \Lambda) \cdot \epsilon \geq \eta$$

$$\epsilon \geq 0$$



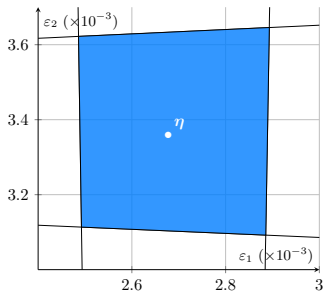
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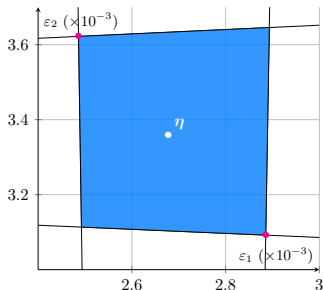
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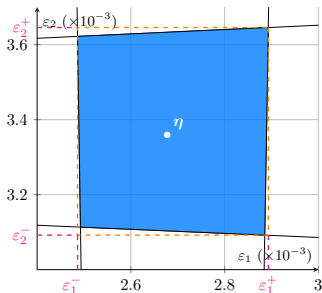
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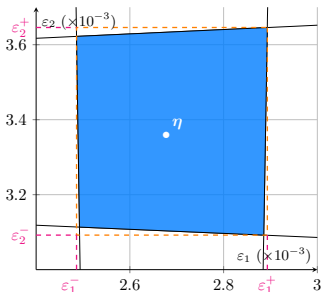
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Lower Bounds for Perov Theorem

For all $i \in \llbracket 1, p \rrbracket$,

$$d(x, x^*)_i = \varepsilon_i \geq \varepsilon_i^-$$

with ε_i^- given by the intersection of the i -th lower bound constraint together with all the j -th upper bound constraints, for $j \neq i$.



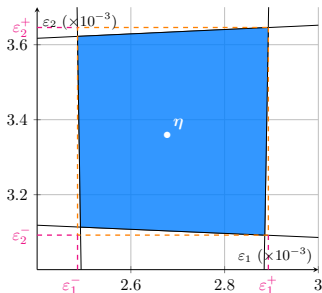
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Example

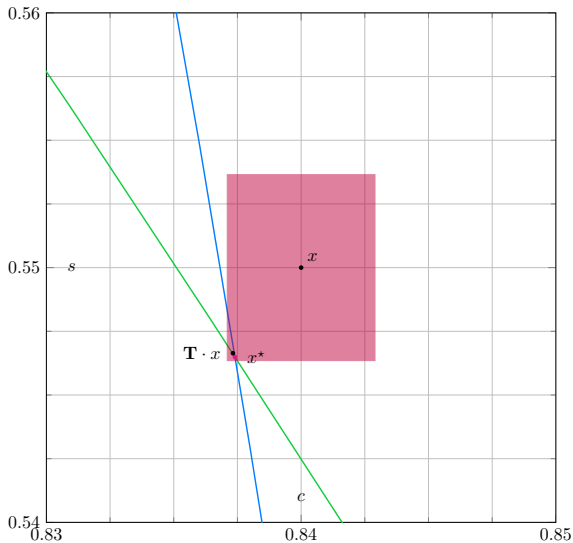
Enclosures obtained by the theorem:

$$\varepsilon_1^- = 2.48 \cdot 10^{-3} \quad \varepsilon_1^+ = 2.90 \cdot 10^{-3}$$

$$\varepsilon_2^- = 3.09 \cdot 10^{-3} \quad \varepsilon_2^+ = 3.65 \cdot 10^{-3}$$

Example: Polynomial Equation in the Plane

Componentwise Tight Error Bounds



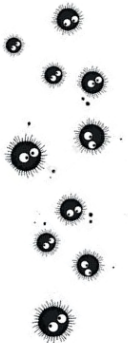
Outline

1 Introduction

2 Multinorm Validation: a New Framework

3 A Posteriori Validation of Vector-Valued D-Finite Functions

4 Conclusion and Future Work



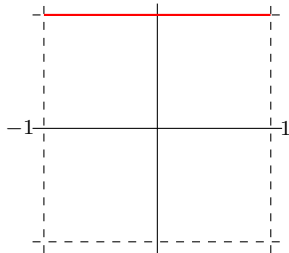


Chebyshev Family of Polynomials

$$T_0(X) = 1,$$

$$T_1(X) = X,$$

$$T_{n+2}(X) = 2XT_{n+1}(X) - T_n(X).$$



$$T_0(X) = 1$$

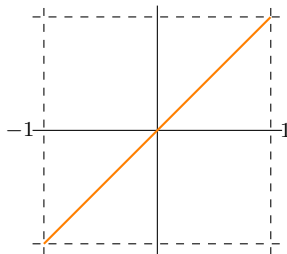


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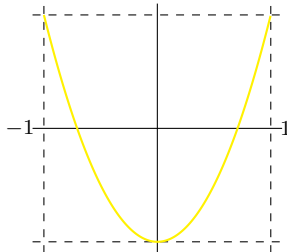


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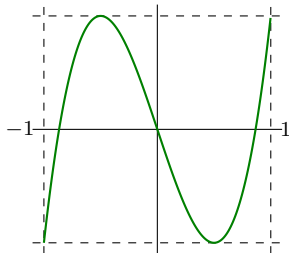


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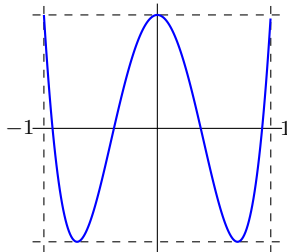


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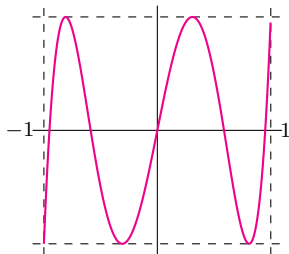


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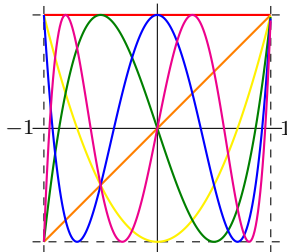
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$$\Rightarrow \forall t \in [-1, 1], |T_n(t)| \leq 1.$$



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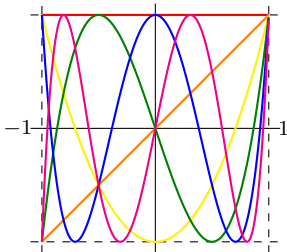
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$$\langle f, g \rangle = \int_{-1}^1 \frac{f(t)g(t)}{\sqrt{1-t^2}} dt = \int_0^\pi f(\cos \vartheta)g(\cos \vartheta) d\vartheta.$$

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$$\blacksquare \widehat{f}^{[N]}(t) = \sum_{n \leq N} a_n T_n(t), \quad t \in [-1, 1].$$



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Convergence Theorems

$$\blacksquare \text{If } f \in \mathcal{C}^k, \widehat{f}^{[N]} \rightarrow f \text{ in } O(N^{-k}).$$

$$\blacksquare \text{If } f \text{ analytic, } \widehat{f}^{[N]} \rightarrow f \text{ exponentially fast.}$$



Vector-Valued D-Finite Equation and Initial Value Problem

$$\begin{aligned} Y^{(r)}(t) + A_{r-1}(t) \cdot Y^{(r-1)}(t) + \dots + A_1(t) \cdot Y'(t) + A_0(t) \cdot Y(t) &= G(t) \\ Y(-1) = v_0 \quad Y'(-1) = v_1 \quad \dots \quad Y^{(r-1)}(-1) = v_{r-1} &\in \mathbb{R}^p \\ t \in [-1, 1] \quad A_i \in \mathbb{R}[t]^{p \times p}, G \in \mathbb{R}[t]^p. & \end{aligned} \quad (D)$$



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 \end{aligned} \tag{D}$$

Integral Equation with Polynomial Kernel

(D) becomes:

$$Y(t) + \int_{-1}^t \begin{pmatrix} K_{11}(t,s) & \dots & K_{1p}(t,s) \\ \vdots & \ddots & \vdots \\ K_{p1}(t,s) & \dots & K_{pp}(t,s) \end{pmatrix} \cdot Y(s) ds = \Psi(t). \tag{I}$$

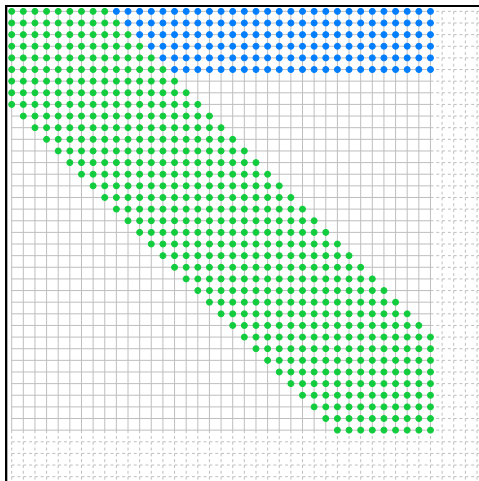
■ $K_{ij} \cdot y(t) = \int_{-1}^t K_{ij}(t,s)y(s)ds$ 1-dimensional *integral operator*.

■ $K = \begin{pmatrix} K_{11} & \dots & K_{1p} \\ \vdots & \ddots & \vdots \\ K_{p1} & \dots & K_{pp} \end{pmatrix}$ p-dimensional *integral operator*.

Compactness and Almost-Banded Structure of \mathbf{K}



$$\mathbf{K}_{ij} \cdot \sum_{k \geq 0} c_k T_k \simeq$$



$$\cdot \begin{pmatrix} c_0 \\ c_1 \\ c_2 \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ c_N \\ c_{N+1} \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \end{pmatrix}$$

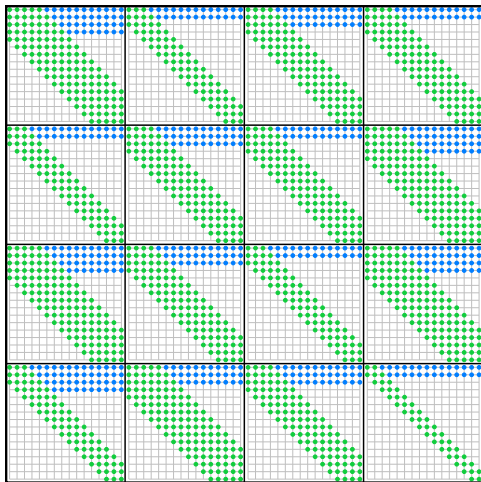
\mathbf{K}_{ij} is **almost-banded** and **compact**.

$$\begin{pmatrix} c_0 \\ c_1 \\ c_2 \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ c_N \\ 0 \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \end{pmatrix}$$

Compactness and Almost-Banded Structure of \mathbf{K}



$$\mathbf{K}^{[N]}. \begin{pmatrix} \sum_{k \geq 0} c_{1k} T_k \\ \vdots \\ \sum_{k \geq 0} c_{pk} T_k \end{pmatrix} \approx$$



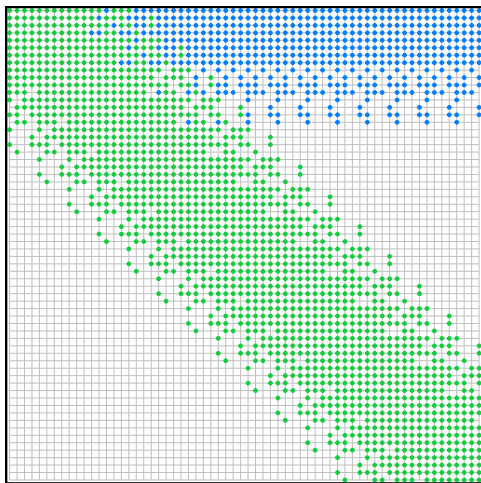
$$\begin{pmatrix} c_{10} \\ c_{11} \\ \vdots \\ c_{1N} \\ \hline \vdots \\ \vdots \\ \vdots \\ \vdots \\ \hline c_{p0} \\ c_{p1} \\ \vdots \\ c_{pN} \end{pmatrix}$$

truncation $\mathbf{K}^{[N]}$ by blocks.

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$$\begin{pmatrix} c_{10} \\ c_{20} \\ \vdots \\ \vdots \\ c_{p0} \\ c_{11} \\ c_{21} \\ \vdots \\ \vdots \\ c_{p1} \\ \vdots \\ \vdots \\ \vdots \\ c_{1N} \\ c_{2N} \\ \vdots \\ \vdots \\ c_{pN} \end{pmatrix}$$

$\mathbf{K}^{[N]}$ in reordered basis.

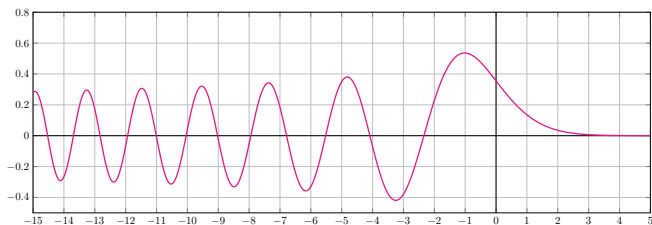
Example: Airy Function

Airy Equation and Integral Reformulation



- Airy function Ai defined by:

$$y'' - ty = 0, \quad \text{Ai}(0) = v_0 \quad \text{and} \quad \text{Ai}'(0) = v_1$$



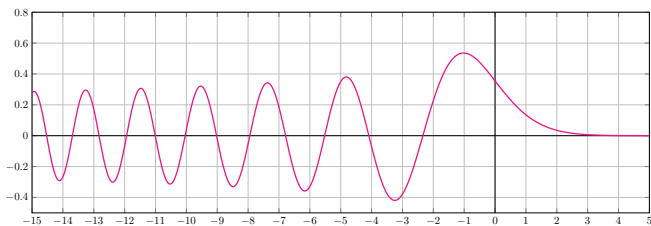
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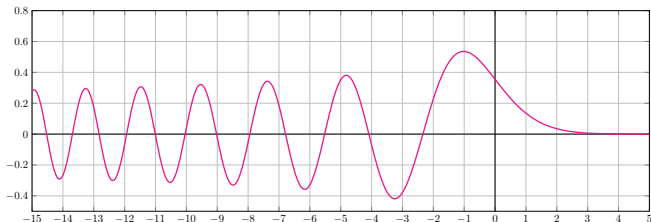
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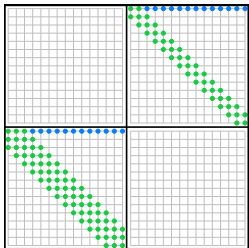
$$Y(t) + \int_{-1}^t \begin{pmatrix} 0 & \frac{2}{3} \\ -\frac{a^2}{4}(s+1) & 0 \end{pmatrix} \cdot Y(s) ds = \begin{pmatrix} v_0 \\ v_1 \end{pmatrix} \Rightarrow Y^*(t) = \begin{pmatrix} \text{Ai}\left(-\frac{a}{2}(t+1)\right) \\ \text{Ai}'\left(-\frac{a}{2}(t+1)\right) \end{pmatrix}.$$

Example: Airy Function

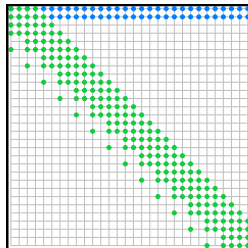
Approximation with Chebyshev Series



- Truncation at order $N = 14$:



Truncated operator $K^{[N]}$ by blocks



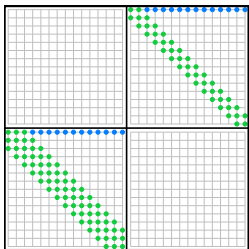
$K^{[N]}$ in reordered basis

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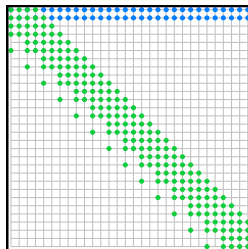
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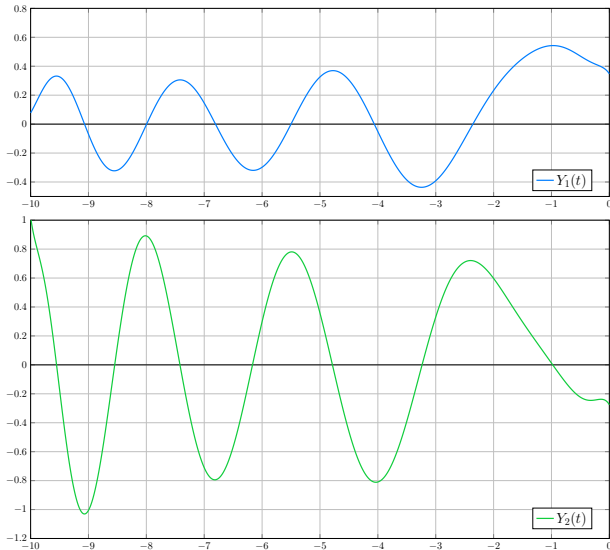
$\mathbf{K}^{[N]}$ in reordered basis

- Obtained approximations for $a = 10$:

$$\begin{aligned} Y_1 = & +0.139 T_0 - 0.152 T_1 + 0.200 T_2 - 0.016 T_3 - 0.010 T_4 + 0.129 T_5 - 0.112 T_6 - 0.032 T_7 \\ & + 0.031 T_8 - 0.162 T_9 - 0.111 T_{10} + 0.103 T_{11} + 0.110 T_{12} - 0.005 T_{13} - 0.033 T_{14} \\ Y_2 = & +0.057 T_0 + 0.130 T_1 + 0.052 T_2 + 0.290 T_3 + 0.033 T_4 + 0.273 T_5 + 0.291 T_6 + 0.004 T_7 \\ & + 0.203 T_8 + 0.104 T_9 - 0.380 T_{10} - 0.340 T_{11} + 0.073 T_{12} + 0.187 T_{13} + 0.044 T_{14} \end{aligned}$$

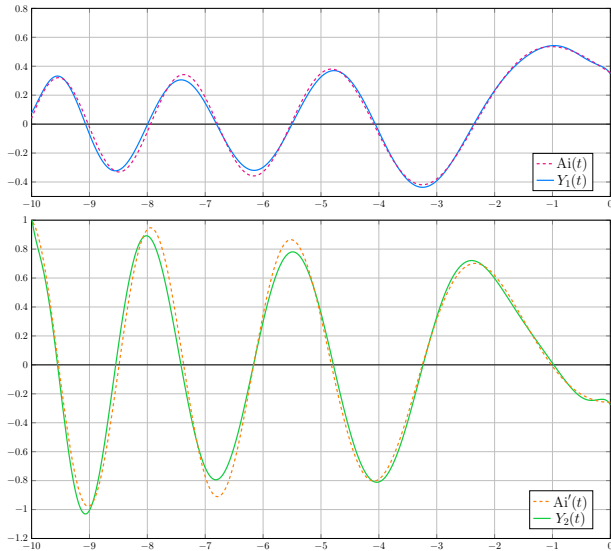
Example: Airy Function

Plots



Example: Airy Function

Plots





Construct \mathbf{T}

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- Approx inverse:

$$\mathbf{A} \approx (\mathbf{I} + \mathbf{K}^{[N_v]})^{-1}$$



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Ψ^1 Banach Space

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Construct \mathbf{T}

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Decomposition of the Operator Norm

$$\begin{aligned} \|\mathbf{DT}\|_{(\Psi^1)^p} &= \|\mathbf{1} - \mathbf{A} \cdot (\mathbf{1} + \mathbf{K})\|_{(\Psi^1)^p} \\ &\leq \underbrace{\|\mathbf{1} - \mathbf{A} \cdot (\mathbf{1} + \mathbf{K}^{[N_v]})\|_{(\Psi^1)^p}}_{\text{Approximation error}} + \underbrace{\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\|_{(\Psi^1)^p}}_{\text{Truncation error}}. \end{aligned}$$

Ψ^1 Banach Space

- $\|y\|_{\Psi^1} = \sum_{n \geq 0} |[y]_n| \geq \|y\|_{\infty}$.
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Construct \mathbf{T}

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Approximation error:

- Finite-dimensional problem.
- Matrix multiplications and Ψ^1 -norm.

Truncation error:

- Infinite-dimensional problem.
- Crude bounds \Rightarrow large N_v .
- Smart bounding techniques.



Rigorous Chebyshev Approximation - Summary

- 1 Integral reformulation,
- 2 Numerical approximation Y of Y^* ,
- 3 Creating Newton-like operator \mathbf{T} ,
- 4 Computing $\Lambda \geq \|D\mathbf{T}\|_{(\mathcal{U}^1)^p}$,
- 5 If $\rho(\Lambda) < 1$, bound $\|Y - \mathbf{T} \cdot T\|_{(\mathcal{U}^1)^p}$ and apply Perov theorem.

Example: Airy Function

Validation with Newton-like Method



Rigorous Chebyshev Approximation - Summary

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Example: Airy Function over $[-10, 0]$

► with $N_v = 1000$:

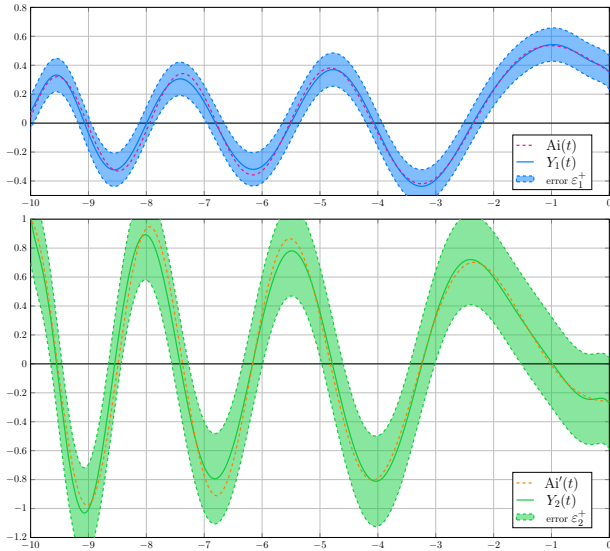
$$\Lambda = \begin{pmatrix} 7.56 \cdot 10^{-4} & 8.71 \cdot 10^{-3} \\ 3.92 \cdot 10^{-2} & 1.11 \cdot 10^{-2} \end{pmatrix}$$

► $\varepsilon_1^- \leq \|Y_1 - \text{Ai}\|_{\mathcal{Y}^1} \leq \varepsilon_1^+$ and
 $\varepsilon_2^- \leq \|Y_2 - \text{Ai}'\|_{\mathcal{Y}^1} \leq \varepsilon_2^+$ with:

$$\begin{array}{ll} \varepsilon_1^- = 0.109 & \varepsilon_1^+ = 0.115 \\ \varepsilon_2^- = 0.296 & \varepsilon_2^+ = 0.312 \end{array}$$

Example: Airy Function

Error Tubes



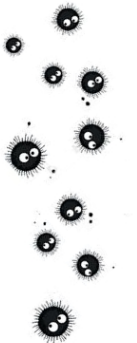
Outline

1 Introduction

2 Multinorm Validation: a New Framework

3 A Posteriori Validation of Vector-Valued D-Finite Functions

4 Conclusion and Future Work





- A general framework for **multinorm validation**.
- An algorithm for **Rigorous Polynomial Approximations** to vector-valued D-finite functions.
- Generalization to **non-polynomial** systems of linear ODEs.
- C library freely available at <https://gforge.inria.fr/projects/tchebyapprox>.
- Towards a **certified Coq** implementation.



Lower Bounds for Perov Theorem

If \mathbf{T} is $\mathbf{\Lambda}$ -Lipschitz with $\mathbf{\Lambda}$ convergent to zero, then for all $i \in \llbracket 1, p \rrbracket$:

$$d(x, x^*)_i \geq \varepsilon_i^- = \left((\mathbf{1} - \mathbf{D}_i \cdot \mathbf{\Lambda})^{-1} \cdot d(x, \mathbf{T} \cdot x) \right)_i \quad \text{with} \quad \mathbf{D}_i = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & \boxed{-1} & \\ & & & \ddots \\ & & & & 1 \end{pmatrix}.$$

Sketch of the proof:

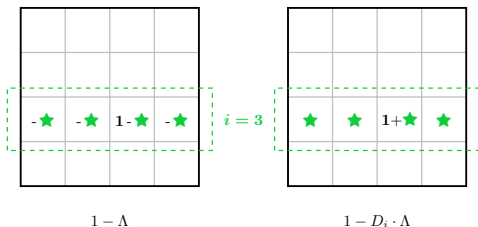


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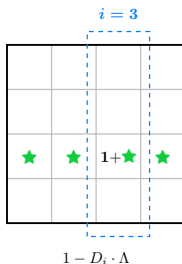
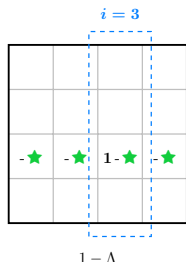


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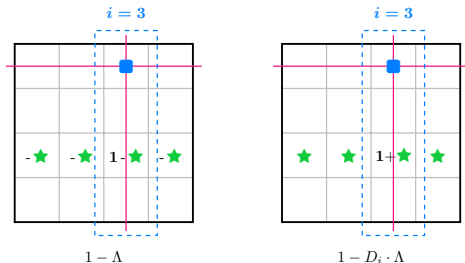


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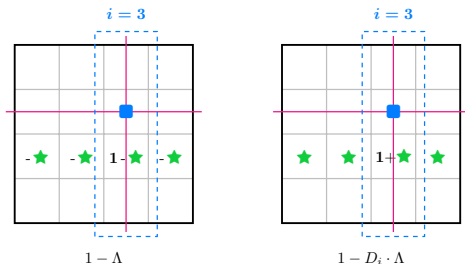


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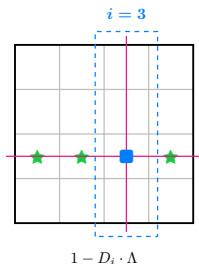
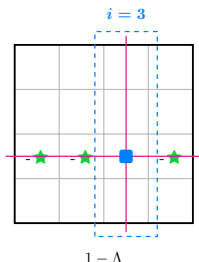


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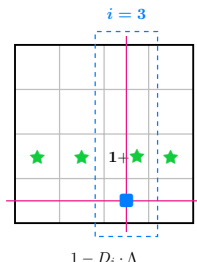
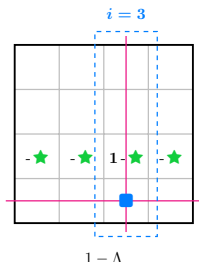


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- Overapproximation ratio:

$$\frac{\varepsilon_i^+}{\varepsilon_i^-} = \frac{d' \, c_i \eta_i + \sum_{j \neq i} c_j \eta_j}{d \, c_i \eta_i - \sum_{j \neq i} c_j \eta_j},$$

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$$\mathcal{C}_\kappa =$$

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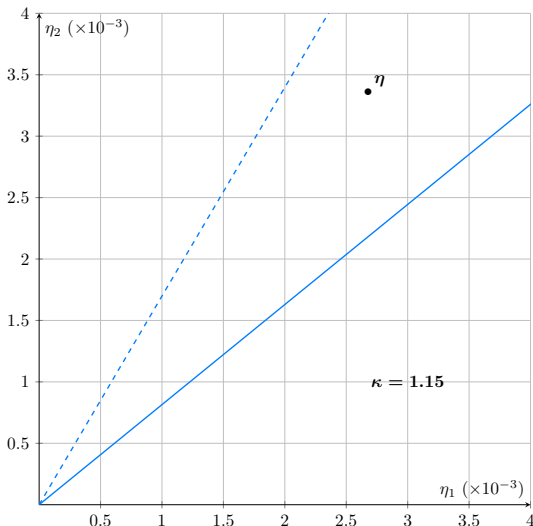
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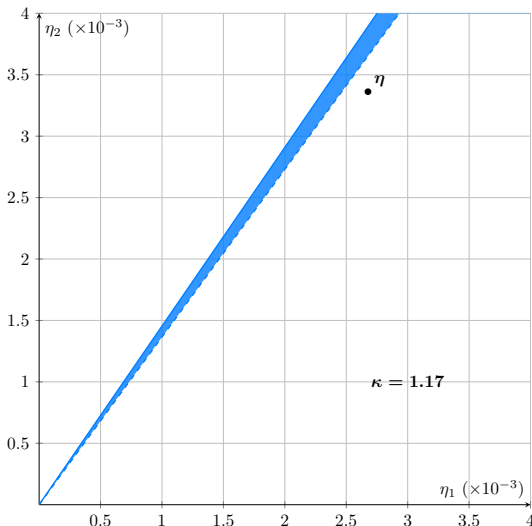
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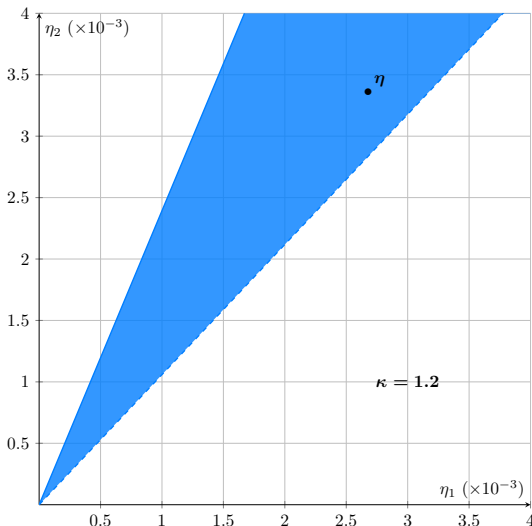
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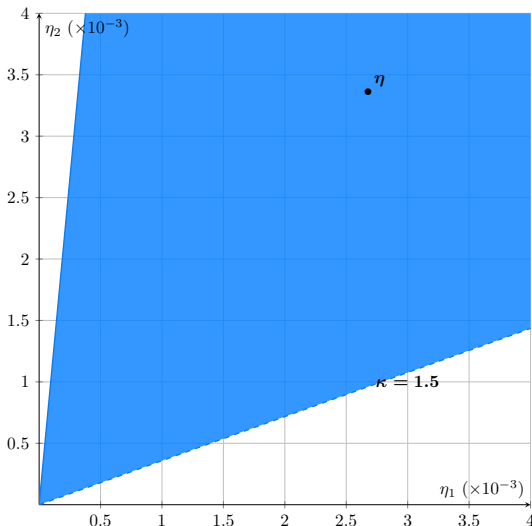
$$d' = \det(1 - D_i \cdot \Lambda).$$

- $\frac{\varepsilon_i^+}{\varepsilon_i^-} \leq \kappa \Leftrightarrow$

$$\eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j$$

Tightness Cone

$$\mathcal{C}_\kappa = \bigcap_{1 \leq i \leq p} \left\{ \eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j \right\}$$





- Overapproximation ratio:

$$\frac{\varepsilon_i^+}{\varepsilon_i^-} = \frac{d' c_i \eta_i + \sum_{j \neq i} c_j \eta_j}{d c_i \eta_i - \sum_{j \neq i} c_j \eta_j},$$

$$c_j = (1 - \Lambda)_{ij}^{-1}, \quad d = \det(1 - \Lambda), \\ d' = \det(1 - D_i \cdot \Lambda).$$

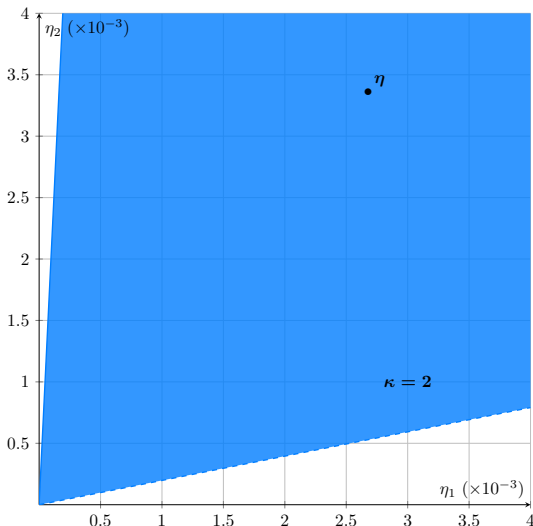
$$\frac{\varepsilon_i^+}{\varepsilon_i^-} \leq \kappa \Leftrightarrow$$

$$\eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j$$

Tightness Cone

$$\mathcal{C}_\kappa =$$

$$\bigcap_{1 \leq i \leq p} \left\{ \eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j \right\}$$





- ▶ Overapproximation ratio:

$$\frac{\varepsilon_i^+}{\varepsilon_i^-} = \frac{d' \ c_i \eta_i + \sum_{j \neq i} c_j \eta_j}{d \ c_i \eta_i - \sum_{j \neq i} c_j \eta_j},$$

$$c_j = (1 - \mathbf{A})_{jj}^{-1}, \quad d = \det(1 - \mathbf{A}),$$

$$d' = \det(1 - \mathbf{D}_i \cdot \mathbf{A}).$$

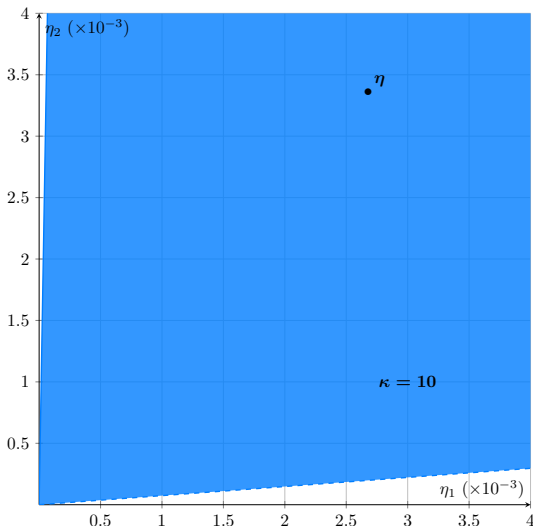
- ▶ $\frac{\varepsilon_i^+}{\varepsilon_i^-} \leq \kappa \Leftrightarrow$

$$\eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j$$

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$$c_j = (1 - \Lambda)_{ij}^{-1}, \quad d = \det(1 - \Lambda), \\ d' = \det(1 - D_i \cdot \Lambda).$$

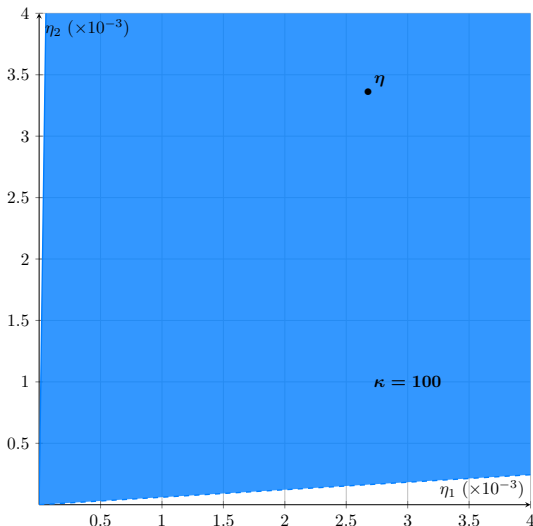
- ▶ $\frac{\varepsilon_i^+}{\varepsilon_i^-} \leq \kappa \Leftrightarrow$

$$\eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j$$

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Tightness of Error Enclosures [Appendix]



- ▶ Overapproximation ratio:

$$\frac{\varepsilon_i^+}{\varepsilon_i^-} = \frac{d' c_i \eta_i + \sum_{j \neq i} c_j \eta_j}{d c_i \eta_i - \sum_{j \neq i} c_j \eta_j},$$

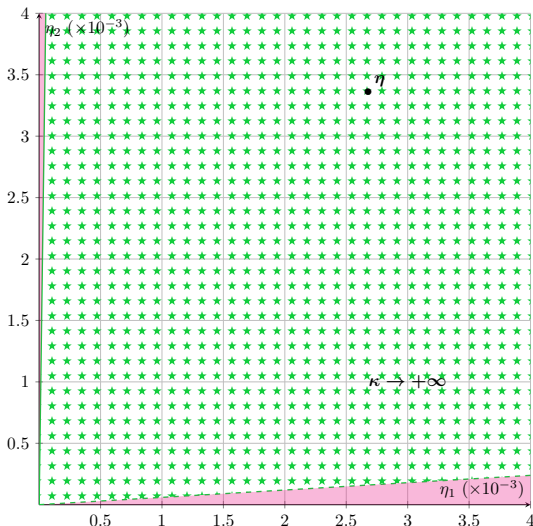
$$c_j = (1 - \Lambda)_{ij}^{-1}, \quad d = \det(1 - \Lambda), \\ d' = \det(1 - D_i \cdot \Lambda).$$

- ▶ $\frac{\varepsilon_i^+}{\varepsilon_i^-} \leq \kappa \Leftrightarrow$

$$\eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j$$

Tightness Cone

$$\mathcal{C}_\kappa = \bigcap_{1 \leq i \leq p} \left\{ \eta_i \geq \frac{\kappa d + d'}{\kappa d - d'} \frac{1}{c_i} \sum_{j \neq i} c_j \eta_j \right\}$$





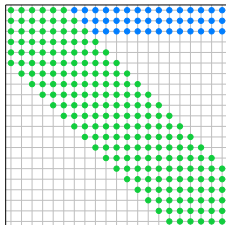
Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$



Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$

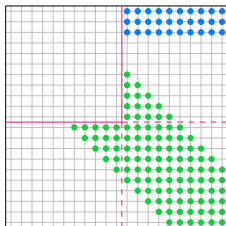


\mathbf{K}



Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$



$\mathbf{K} - \mathbf{K}^{[N]}$

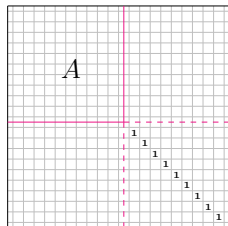
Designing the Newton-like Operator \mathbf{T}

Bounding the Truncation Error

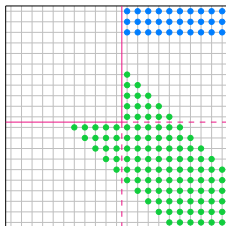


Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$



\mathbf{A}



$\mathbf{K} - \mathbf{K}^{[N]}$

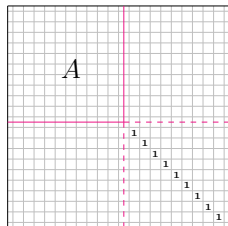
Designing the Newton-like Operator \mathbf{T}

Bounding the Truncation Error

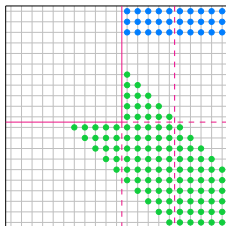


Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$



A



$\mathbf{K} - \mathbf{K}^{[N]}$

Designing the Newton-like Operator \mathbf{T}

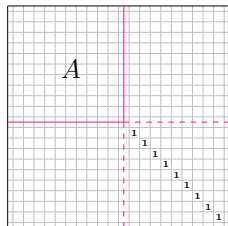
Bounding the Truncation Error



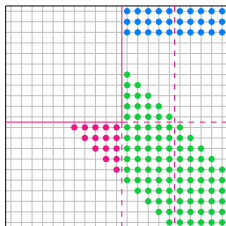
Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$

1 Direct computation.



A



$\mathbf{K} - \mathbf{K}^{[N]}$

Designing the Newton-like Operator \mathbf{T}

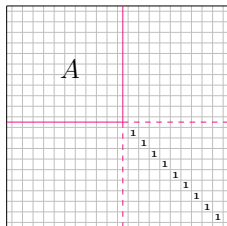
Bounding the Truncation Error



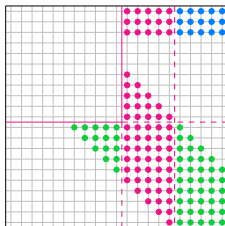
Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$

- 1 Direct computation.
- 2 Direct computation.



\mathbf{A}



$\mathbf{K} - \mathbf{K}^{[N]}$

Designing the Newton-like Operator \mathbf{T}

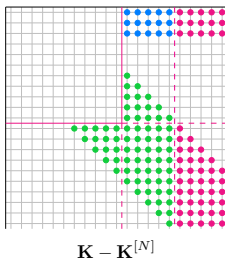
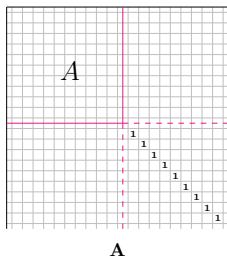
Bounding the Truncation Error



Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$

- 1 Direct computation.
- 2 Direct computation.
- 3 Bound the remaining *infinite* number of columns:





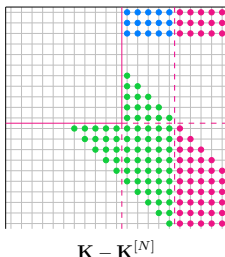
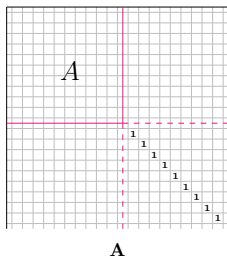
Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$

- 1 Direct computation.
- 2 Direct computation.
- 3 Bound the remaining *infinite* number of columns:

- Using the bounds in $1/i$ and $1/i^2$: possibly large overestimations.

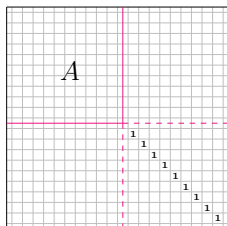
$$\text{diag}(i) \leq \frac{C}{i} \quad \text{init}(i) \leq \frac{D}{i^2}$$



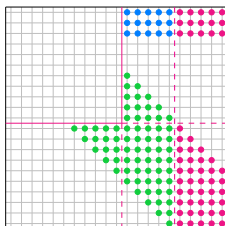


Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$



\mathbf{A}



$\mathbf{K} - \mathbf{K}^{[N]}$

- 1 Direct computation.
- 2 Direct computation.
- 3 Bound the remaining *infinite* number of columns:

- Using the bounds in $1/i$ and $1/i^2$: possibly large overestimations.

$$\text{diag}(i) \leq \frac{C}{i} \quad \text{init}(i) \leq \frac{D}{i^2}$$

- Using a first order difference method: differences in $1/i^2$ and $1/i^4$.

$$\text{diag}(i) \leq \text{diag}(i_0) + \frac{C'}{i^2}$$

$$\text{init}(i) \leq \text{init}(i_0) + \frac{D'}{i^4}$$

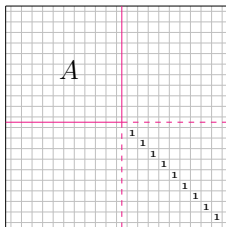
Designing the Newton-like Operator \mathbf{T}

Bounding the Truncation Error

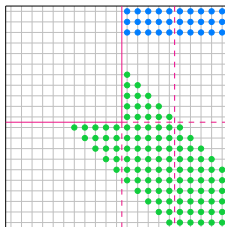


Truncation Error

$$\|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]})\| = \sup_{i \geq 0} \|\mathbf{A} \cdot (\mathbf{K} - \mathbf{K}^{[N_v]}) \cdot \mathbf{T}_i\|$$



\mathbf{A}



$\mathbf{K} - \mathbf{K}^{[N]}$

- 1 Direct computation.
- 2 Direct computation.
- 3 Bound the remaining *infinite* number of columns:
 - Using the bounds in $1/i$ and $1/i^2$: possibly large overestimations.

$$\text{diag}(i) \leq \frac{C}{i} \quad \text{init}(i) \leq \frac{D}{i^2}$$

- Using a first order difference method: differences in $1/i^2$ and $1/i^4$.

$$\text{diag}(i) \leq \text{diag}(i_0) + \frac{C'}{i^2}$$

$$\text{init}(i) \leq \text{init}(i_0) + \frac{D'}{i^4}$$